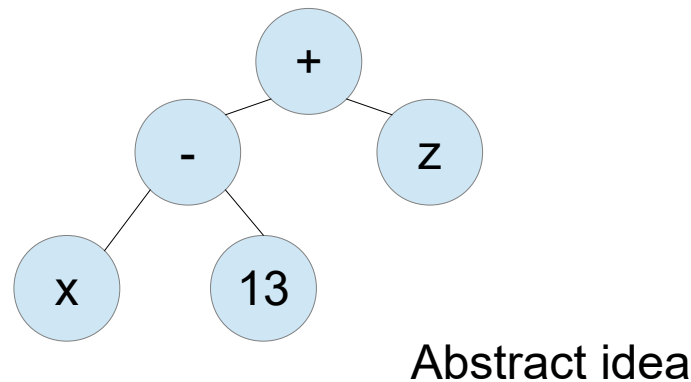
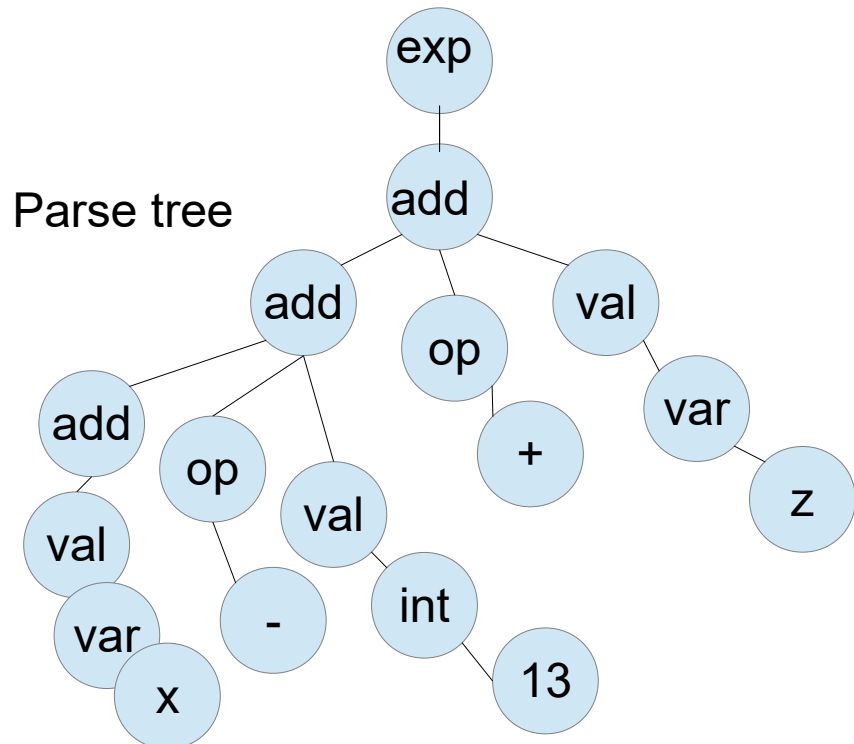


IR trees/graphs

- Various possible tree/graph intermediate representations:
- Parse tree: directly based on grammar of the source language
- Syntax tree: abstract from parse tree (less language dependent)
- Dependency graph: show hierarchy of declared/defined items, and which ones depend on which others
- Control flow graphs: divide code into uninterrupted blocks of code, with directed edges indicating possible flow between them
- Call graphs: nodes for procedures, directed edges indicate calls

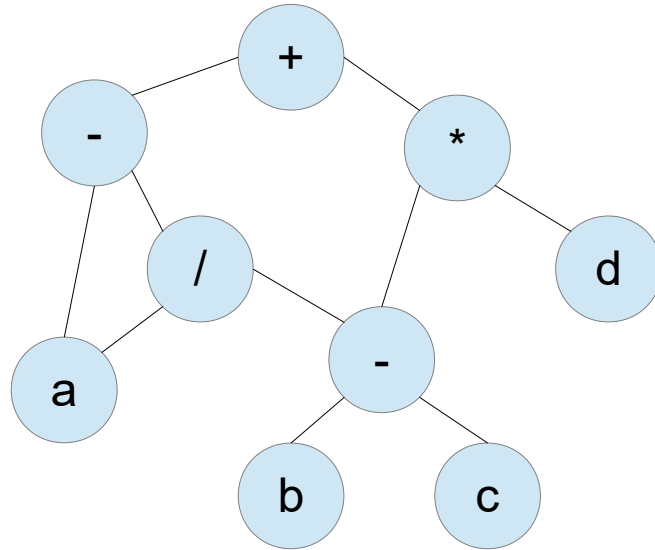
Parse trees vs syntax trees

- Attempt to abstract data and operations away from language specific grammar rules



Directed acyclic syntax graphs

- Reduce size by identifying/reusing common subtrees
- $(a - (a / (b - c))) + ((b - c) * d)$



Syntax DAGs

- Need an effective way to recognize when common subtrees exist, ideally asap during construction
- Introduces many possible language-independent optimizations, e.g.
 - in expressions: store result of subtree in a temp variable rather than recomputing
 - if subtree represents a block of statements then replace with callable function)

Tree construction from grammar

- Suppose we build our tree with two kinds of nodes:
 - leaf: holds the tokens in our source grammar
 - node: internal node, corresponding to nonterminals
- For each grammar rule, we define a rule on how to build the appropriate tree node/leaf
- For top-down derivation, we start with a node for our top level nonterminal, and on each rule application we apply the appropriate construction

Example: construction rules

<code>expr-->addx</code>	<code>expr.node = addx.node</code>
<code>addx--> addx aop multx</code>	<code>addx.node = new node(aop.node, addx.node, multx.node)</code>
<code>addx --> multx</code>	<code>addx.node = multx.node</code>
<code>multx --> multx mop valx</code>	<code>multx.node = new node(mop.node, multx.node, valx.node)</code>
<code>multx --> valx</code>	<code>multx.node = valx.node</code>
<code>valx --> VAR NUM</code>	<code>valx.node = new leaf (VAR, VAR.txt)</code> <code>valx.node = new leaf (NUM, NUM.val)</code>
<code>valx --> '(' expr ')'</code>	<code>valx.node = new node('(', expr.node, ')')</code>
<code>aop --> '+' '-'</code>	<code>aop.node = new leaf('+')</code> <code>aop.node = new leaf('-')</code>
<code>mop --> '*' '/'</code>	<code>mop.node = new leaf('*')</code> <code>mop.node = new leaf('/')</code>

Array-of-records implementation

- Need a way to represent our leaf/node collection, e.g.
 - leaf record type
 - node record type
 - keep an array of records (and counter)
- Each leaf/node thus has a unique index value (array pos)
- Cross references between nodes can use the index (giving small storage, fast lookups)
- Often referred to as value-number method, each node has unique associated index number

Value-number example

$i = i + x * 10$

i.e.

$i = (i + (x * 10))$

index	Node/leaf	data1	data2
0	(leaf x)	symtable ptr for x	
1	(leaf i)	symtable ptr for i	
2	(leaf 10)	literal 10	
3	(node *)	0 (index of node x)	2 (index of node 10)
4	(node +)	1 (index of node i)	3 (index of node *)
5	(node =)	1 (index of node i)	4 (index of node +)

Searching problem:

- As we're building the array, we need to search current array content to find operand indices, e.g. to fill in fields for $x * 10$ we need to find indices for node x and 10
- Currently that means a linear search: $O(n)$
- Could store the nodes as a binary search tree instead of an array, so $O(\log(n))$
- Could store the nodes in hash table: collection of buckets, with hash function mapping the operands (e.g. $*$, X , 10) to a bucket, then just linear search the bucket if not empty

Using duplicate subtrees (DAG)

- When building an entry, and have searched for the correct operand indices, look at fields for new entry, check if there's already a matching entry
- e.g. Suppose we have a new entry using $x*10$ again, we look for a (node $*$) with data fields 0 and 2, and find index 3 already provides it

index	Node/leaf	data1	data2
0	(leaf x)	symtable ptr for x	
1	(leaf i)	symtable ptr for i	
2	(leaf 10)	literal 10	
3	(node $*$)	0 (index of node x)	2 (index of node 10)
4	(node $+$)	1 (index of node i)	3 (index of node $*$)
5	(node $=$)	1 (index of node i)	4 (index of node $+$)