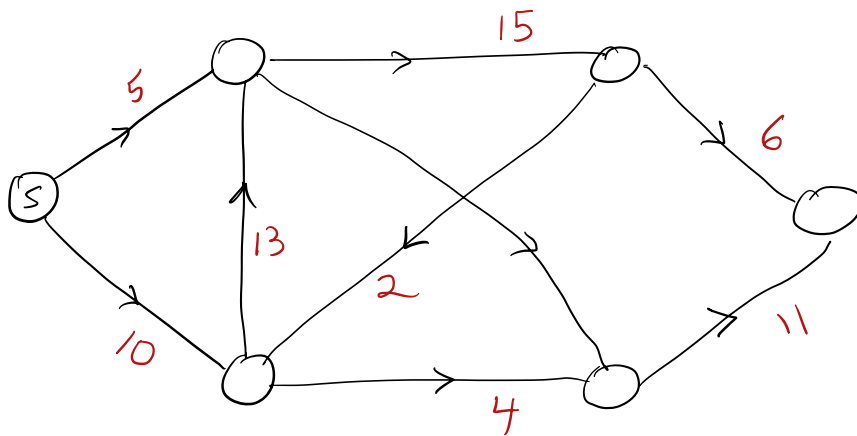


Network Flows

10.25

Given: A directed graph $G=(V,E)$
with edge capacities $C: E \rightarrow \mathbb{R}$
and distinct source vertex (no in-edges) s
and distinct sink vertex (no out-edges) t

Find: Max amount of flow from s to t

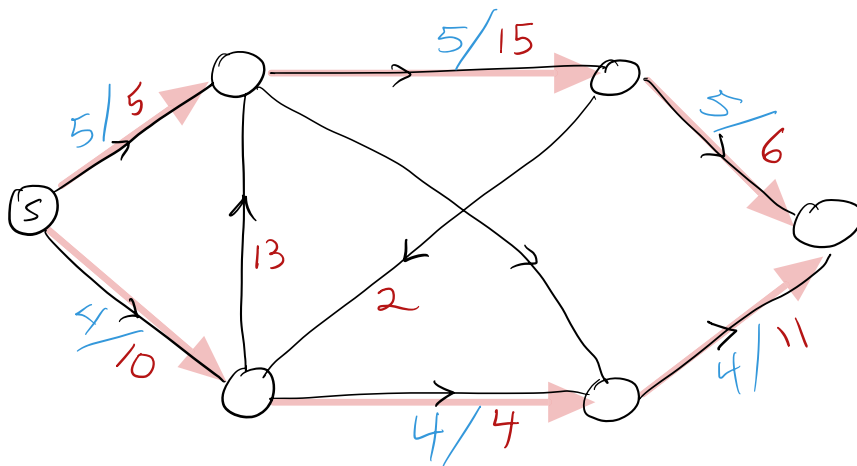


Network Flows

10.25

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Find: Max amount of flow from s to t



$$\begin{aligned} \text{Total flow} &= \text{flow out from } s \\ &= \text{flow into } t \\ &= 9 \end{aligned}$$

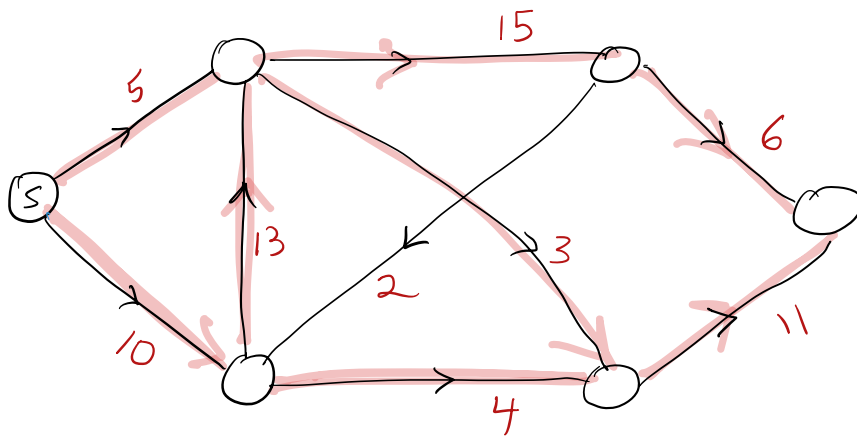
Max?

Network Flows

10.25

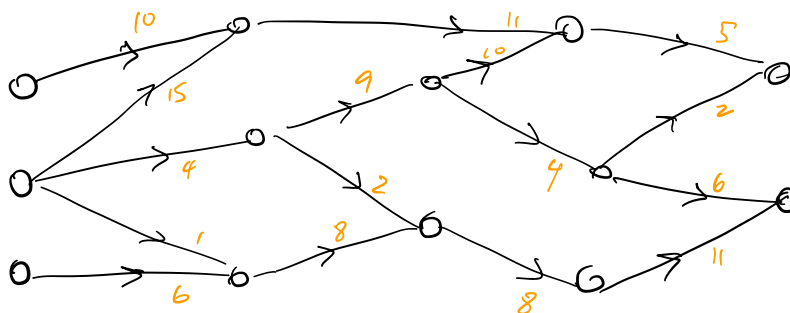
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and distinct source vertex (no in-edges) s
and distinct sink vertex (no out-edges) t

Find: Max amount of flow from s to t



Multi-Source / Multi sink ?

Not a problem.



Defn Flow

A flow is an assignment of values $f: E \rightarrow \mathbb{R}$

such that:

$$\forall (u,v) \in E, f(u,v) \leq c(u,v)$$

Obey
Capacity
Constraints

$$\forall v \in V, v \neq s, v \neq t$$

$$\sum_{u \in V} f(u,v) = \sum_{w \in V} f(v,w)$$

Conservation
of
Flow

Observe: max flow consists in maximizing

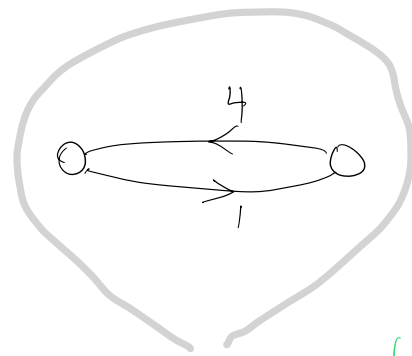
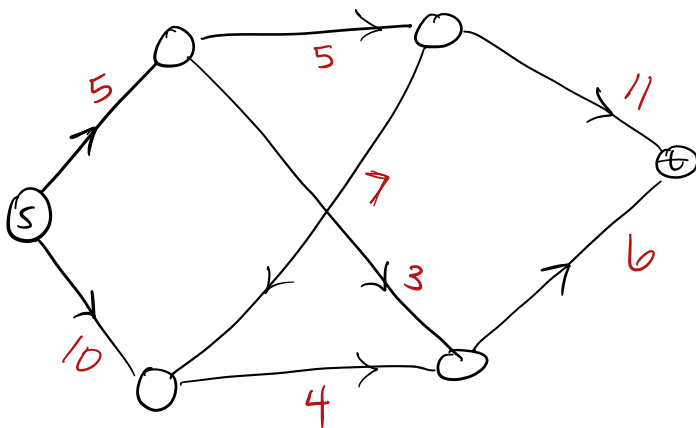
$$\sum_{u \in V} f(s,u)$$

(Why not maximize $\sum_{w \in V} f(w,t)$?

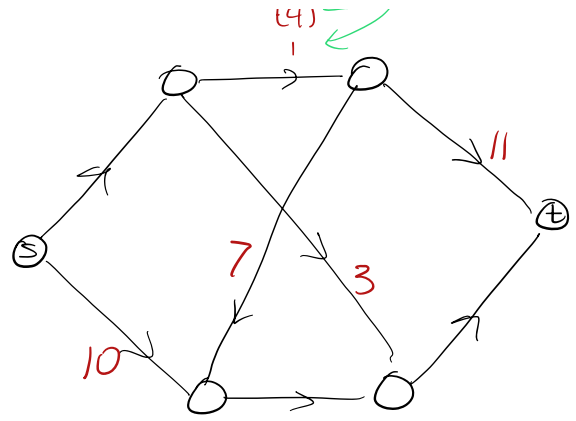
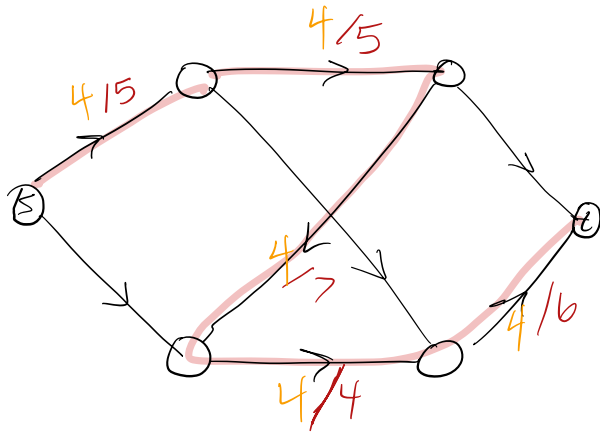
"Flow Augmenting Path" Alg : Ford & Fulkerson

Iterative.

- Find a s - t path that has some **unused capacity** on all the forward edges, and positive flow on all the backwards edges
- Augment current flow with min value unused capacity (or positive flow on backwards edges) on all path edges.



flow (current)
unused capacity



Residual

