

Recall that for a string w , $\#_\sigma(w)$ is the number of occurrences of σ in w .

1. (2 mark) Either find, if they exist, a regular language L_1 and a context-free language L_2 whose intersection is regular, or explain why no such two languages exist.

2. (4 marks) True or False:
 - (a) ___ The union of a (possibly infinite) number of regular languages can be non-regular
 - (b) ___ If L_1 is regular and L_2 is context-free, then $L_1 \cup L_2$ is necessarily also context-free.
 - (c) ___ Given a regular grammar G , each string in $L(G)$ has a unique derivation in G .
 - (d) ___ If L is context-free, then so is L^* .

3. (6 marks) Give an algorithm to determine if a given Context-Free Grammar G is *usable*. Recall that a grammar is usable if there exists at least one string of terminals that can be derived from the start symbol of the grammar. In your answer, let the grammar be $G = (V, \Sigma, R, S)$, where
 S is the start symbol,
 Σ is the alphabet of terminals,
 V is set of variables, and
 R is the sets of production rules.
You may refer to the *left side* of a rule, which consists of a single variable, and the *right side* of a rule, which consists of a string from $(V \cup \Sigma)^*$.

4. $L = \{w \in \{a, b\}^* : \text{the latter half of the string contains only } b\text{'s}\}$. An example of a string in the language is $abaababbbbbbb$. In other words, the last a appears in the first half of the string. In odd strings, the middle symbol must be b .

(a) (4 marks) Give a natural PDA for the language given in part (a) – i.e., not a bottom-up or top-down parser obtained from your grammar.

(b) (4 marks) Give a Context-Free Grammar for the language given above.

5. (4 marks) Give a Context-Free Grammar for the following: $\{a^{i+j}b^i c^j : i, j \geq 0\}$

6. (4 marks) Give a natural PDA for the above language – i.e., not a bottom-up or top-down parser obtained from your grammar.

7. (4 marks) Give a Context-Free Grammar for the following: $\{a^n b^m : 2n = 5m - 3, \text{ and } n, m \geq 0\}$

8. (4 marks) Give a Context-Free Grammar for the following: $\{w \in \{a, b\}^* : w = x a^n b^n x^R, \text{ where } x \text{ is a string in } \{a, b\}^*\}$. Examples of such strings are: ab (where $x = \epsilon = x^R$), and $aababbaaabbabbabaa$ (where $x = aababba$ and $x^R = abbabaa$).

9. (4 marks) Give a natural PDA for the language in the previous question – i.e., not a bottom-up or top-down parser obtained from your grammar.

10. (4 marks) Give a Context-Free Grammar for the following: $\{w \in \{(,)\}^* : \#_l(w) = \#_r(w)$ and every *suffix* of w has at least as many '('s as ')'s}. A suffix is the opposite of a prefix: for a string w , if w can be written as the concatenation of two strings, i.e., $w = w_1w_2$, then the latter part, w_2 is a suffix of w . Note that ϵ is a suffix of every string, and w is a suffix of itself.

11. (4 marks) Give a bottom-up (Shift-Reduce) parser PDA for the following grammar.

$$S \rightarrow X|Y$$

$$X \rightarrow Xc|A$$

$$A \rightarrow aAb|\epsilon$$

$$Y \rightarrow aY|B$$

$$B \rightarrow bBc|\epsilon$$

12. **Bonus:** (3 bonus marks) Prove using the Pumping Lemma that the language L from question 6 is not regular; recall
 $L = \{w \in \{a, b\}^* : \text{the latter half of the string contains only } b\text{'s}\}$. An example of a string in the language is $abaababbbbbbb$. In other words, the last a appears in the first half of the string. In odd strings, the middle symbol must be b .