

From chapter 0...

Sets are like this:  $\{a, b\}$

Annotations:  
 - curly bracket (points to  $\{$  and  $\}$ )  
 - does not imply an order (points to  $a, b$ )

If there is an (implied) "universe"  $U$ ,  
 or ground set from which the set is  
 drawn, then set complementation is defined

$$\overline{S} = U \setminus S$$

Annotations:  
 -  $\overline{S}$ : all elements not in  $S$   
 -  $\setminus$ : "set-minus"  
 -  $S$ : " $S$  complement"

We also have  $\cap$ ,  $\cup$

Sequences and tuples - order is implied

- use  $( \quad )$

Annotations:  
 -  $($ : "open paren"  
 -  $)$ : "close paren"

(Note: graph theory, undirected graphs, abuse this notation somewhat, using  $(u, v)$  for an

undirected edge.)

Review: - Functions + relations  
- Graphs.

## STRINGS & LANGUAGES

An alphabet is a finite, non-empty set.  
Its members are called symbols (or letters).

$\Sigma$  "capital sigma" } often used to  
 $\Gamma$  "capital gamma" } denote alphabets

A string over an alphabet  $\Sigma$  is a  
finite **sequence** of symbols from  $\Sigma$ ,  
duplicates allowed.

$\epsilon$  ("epsilon") is the symbol for the empty  
string.

$$ab \cdot bb = abbb$$

$$S_1 \cdot S_2 = S_1 S_2$$

Juxtaposing two strings is an operation called  
"concatenation", yielding a new string. Also use •

## Recursive Definitions.

Alternative definition of strings over  $\{a, b\}$ .

- Defn:
1.  $\epsilon$  is a string over  $\{a, b\}$
  2. if  $s$  is a string over  $\{a, b\}$ ,  
so is  $s \cdot a$
  3. if  $s$  is a string over  $\{a, b\}$ ,  
so is  $s \cdot b$
  4. Nothing else is a string over  $\{a, b\}$   
 $aab$  a string over  $\{a, b\}$ ?

The above is a recursive definition of strings (over the given alphabet).

Once we have a recursive definition of an object, we can easily define functions of that object.

Def<sup>n</sup> length of a string over  $\{a, b\}$ ,  
denoted with abs-value  $|s|$

1.  $|\epsilon| = 0$

2.  $|s \cdot a| = |s| + 1$ , where  $s$  is a string over  $\{a, b\}$

3.  $|s \cdot b| = |s| + 1$  " " "

Length is now defined for all strings over  $\{a, b\}$ , since (4) "Nothing else is a string over  $\{a, b\}$ "

Def<sup>n</sup> string over  $\Sigma$ .

1.  $\epsilon$  is a string over  $\Sigma$
2. if  $s$  is a string over  $\Sigma$  and  $\sigma \in \Sigma$   
then  $s \cdot \underline{\sigma}$  is a string over  $\Sigma$
3. Nothing else is a string over  $\Sigma$ .

For you to do:

1. Give a definition of string length for strings over  $\Sigma$ .
2. We want  $\#_a(s)$  defined on strings over  $\{a, b\}$  to be the number of  $a$ 's in  $s$ .  
Make it so, using the recursive definition.
3. Do same for  $\#_\sigma(s)$  for strings over  $\Sigma$ ,  $\sigma \in \Sigma$ .
4. (challenge) give a coherent recursive definition of concat  $(\cdot)$  for strings over  $\Sigma$
5. Problem 0.12

Claim: For any set  $S$  of horses,  $|S| = \underline{h} \geq 1$ ,  
the horses in the set are the same

colour. (monochromatic)

"Proof": By induction on  $h$ .

✓  $\rightarrow$  Basis:  $h=1$ . Obvious.

Induction Step:  $\forall$   $|S| = k > 1$

Ind Hyp:  $\forall k' < k$ , any set of  $k'$  horses is monochromatic.

Take one horse <sup>first hor</sup> out of  $S$ , yielding  $S'$ .

$|S'| < k$  so  $S'$  is monochromatic, by IH.

Put that horse back in and take out a different horse <sup>horse 2</sup>, yielding  $S''$ .

$S''$  horses are also all same colour as first horse, by Ind Hyp.

$\therefore$  all horses in  $S$  are same colour. 

What is wrong?



Def<sup>n</sup> of string length for strings over  $\Sigma$ .

1.  $|\varepsilon| = 0$

2.  $|s \cdot \sigma| = |s| + 1 \quad \forall \text{ strings } s$   
over  $\Sigma$ ,  $\forall \sigma \in \Sigma$ .

Def<sup>n</sup> of  $\#_a(s)$ ,  $s$  a string over  $\{a, b\}$

1.  $\#_a(\varepsilon) = 0$

2.  $\#_a(s \cdot a) = \#_a(s) + 1$

3.  $\#_a(s \cdot b) = \#_a(s)$

Def<sup>n</sup> of  $\#_\sigma$ , a function of strings over  $\Sigma$ .  
 $\sigma \in \Sigma$ .

1.  $\#_\sigma(\varepsilon) = 0$

2.  $\#_\sigma(s \cdot \sigma) = \#_\sigma(s) + 1$

3.  $\#_\sigma(s \cdot \delta) = \#_\sigma(s)$

$\delta \neq \sigma$ ,  
 $\delta \in \Sigma$ .