

Jan 28, 2025 Pumping Lemma

- using it to prove languages are not regular -

For each of the following languages, either prove it is regular or prove it is not regular

$$L_1 = \{a^i b^j \mid i, j \geq 0 \text{ and } i+j=5\} \quad \begin{array}{l} bbbbb + abbbb + aabbb + \\ aaabb + aaaab + aaaaa \end{array}$$

$$L_2 = \{a^i b^j \mid i, j \geq 0 \text{ and } i-j=5\}$$

$$L_3 = \{a^i b^j \mid i, j \geq 0 \text{ and } i-j \equiv 0 \pmod{5}\}$$

$$L_4 = \{a^i b^j \mid 0 \leq i \leq j \leq 2000\}$$

$$L_5 = \{w \in \{a, b\}^* \mid w = w_1 w_2 \text{ where } |w_1| = |w_2| \text{ and } \#_a(w_1) = \#_a(w_2)\} \quad \text{Tricky!}$$

$$L_6 = \{w \in \{a-z\}^* \mid \text{every letter in } w \text{ appears at least twice}\}$$

e.g. unprosperousness

$$L_7 = \{w \in \{a, b\}^* \mid \#(w_a) \neq \#(w_b)\}$$

hint: use a closure theorem \Leftarrow we'll cover it today in class

$$L_8 = \{ a^i b^j c^k \mid k = i+j \} \text{ // and } i, j, k \geq 0$$

PL: $\forall RL \ L \ \exists$ a constant p such that $\forall w \in L$,
 $|w| \geq p$, we have: $w = xyz$ such that ⁱ⁾ $|y| > 0$,
ⁱⁱ⁾ $|xy| \leq p$, and $xy^i z \in L \ \forall i = 0, 1, 2, \dots$

aaabbbcccc

Claim:

9. $\{ a^i b^j c^k \mid \text{if } i=1 \text{ then } j=k \}$ Tricky!

$$L_7 = \{ w \in \{a,b\}^* \mid \#_a(w) \neq \#_b(w) \}$$

Claim: L_7 is not regular.

Proof: BWOC. Suppose L_7 is regular.

Then so is $\overline{L_7} \cap L(a^*b^*)$

Observe that $\overline{L_7} \cap L(a^*b^*) = A^n B^n$

But $A^n B^n$ is not regular!

$\Rightarrow \Leftarrow$

$\therefore L_7$ is not regular. \square

$$L_5 = \{ w \in \{a,b\}^* \mid w = w_1 w_2, |w_1| = |w_2|, \text{ and } \#_a(w_1) = \#_a(w_2) \}$$

Claim: L_5 is not regular.

Proof:

