

Recall:

A_{TM} is undecidable.

Proof: ^{BWOC.} $\nexists A_{TM}$ is decided by a TM A

Then we can construct the following TM S

$S =$ " on input $\langle M \rangle$:
1. Run A on input $\langle M, \langle M \rangle \rangle$
- if A accepts, ACCEPT.
- if A rejects, REJECT. "

decider for known undecidable (pointing to S)
decider for language we are trying to prove is undecidable. (pointing to A)

Clearly S accepts $\langle M \rangle$ iff $M \in \text{SelfAcc}$,
and rejects otherwise.

$\circ \circ S$ decides SelfAcc . But SelfAcc is undecidable.

$\Rightarrow \Leftarrow$

$\circ \circ \nexists$ a TM that decides A_{TM} , and so A_{TM} is undecidable. \square

$\text{Halt} = \{ \langle M, w \rangle \mid M \text{ is a TM that halts on input } w \}$

Prove that Halt is undecidable.

Theorem: Halt is undecidable.

Proof: Bwoc. \S Halt is decided by TM H .

Then we can construct the following TM:

$A =$ "on input $\langle M, w \rangle$ "

1. Run H on $\langle M, w \rangle$.

if H rejects, REJECT.

2. if H accepts,

Run M on w .

if M accepts,

if M rejects,

A decides A_{TM} . But A_{TM} is

undecidable. $\Rightarrow \Leftarrow$

$\therefore H$ cannot exist, and

Halt is undecidable.

