

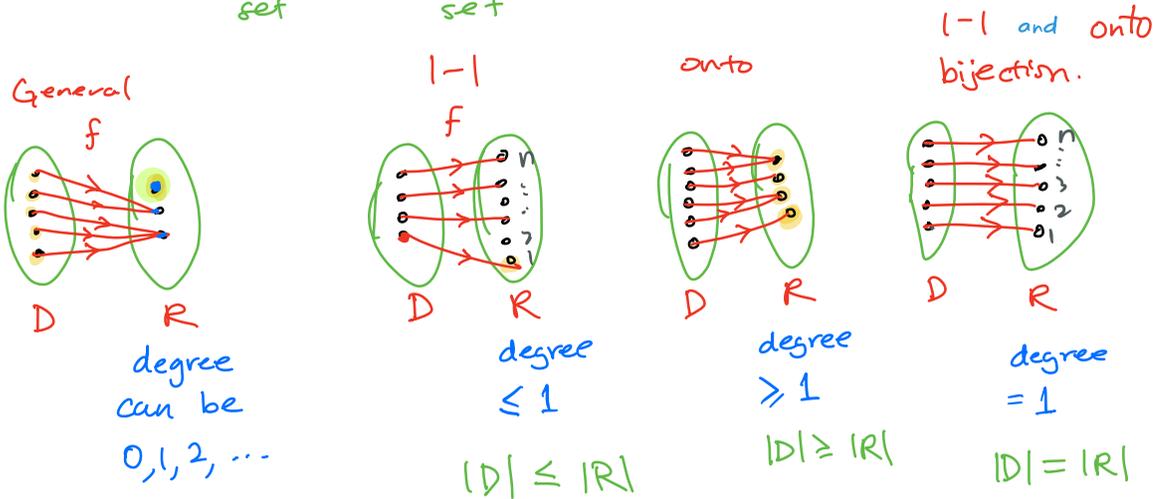
# A short section on counting (Section 4.2 of text)

How do we capture the size of sets (cardinality)?

- when the set is finite
- when the set is infinite?

Recap about functions:  $| - |$   
 + onto      injective  
                  surjective  
 = correspondence      bijective

function: Domain  $\rightarrow$  Range  
                  set                    set  
 $f: \mathbb{N} \rightarrow \mathbb{R}$



Defn A set  $S$  is finite if  $\exists$  integer  $n$  s.t.  
 $\exists$  an 1-1 function from  $S$  to  $\{1, 2, \dots, n\}$   $\leftarrow [n]$

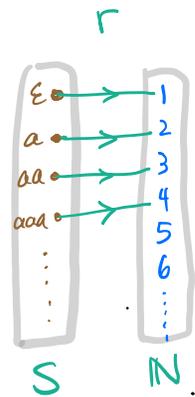
Defn If  $\exists$  a bijection from  $S$  to  $[n]$ , then  
 $|S| = n$  is the cardinality of  $S$

Defn A set is infinite if it is not finite.

Eg.  $\mathbb{N} = \{1, 2, 3, \dots\}$  is infinite set.

Def<sup>n</sup> A set  $S$  is countably infinite if  $\exists$  a bijection  $r$  from  $S$  to  $\mathbb{N}$ . (alternatively  $\mathbb{W}$ )

(In other words,  $S$  is countably infinite if we can enumerate its elements via a ranking function  $r$ .)



$$\text{rank}(aaa) = 4$$

$$r(aaa) = 4$$

$$\text{unrank}(4) = aaa$$

$$r^{-1}(aaa) = 4$$

In an enumeration of  $S$ , every  $x \in S$  has finite rank.

Def<sup>n</sup> A set is countable if it is either finite or countably infinite.

Def<sup>n</sup> An enumeration of  $S$  is a bijection  $r: S \rightarrow \mathbb{N}$ , or  $r: S \rightarrow [n]$  for some  $n$ .

Def<sup>n</sup> A loose enumeration is a 1-1 function  $r: S \rightarrow \mathbb{N}$  or  $[n]$

Eg enumeration of  $L((aa)^*)$

loose enumeration

← tighten

0.  $\epsilon$   
 1.  
 2.  $aa$   
 3.  
 4.  $aaaa$

Each loose enumeration can be "tightened" into an enumeration that starts at 1.

Claim:  $L(a^*)$  is countably infinite.

Proof: 0.  $\epsilon$  in general, string  $a^i$  has  
 1. a rank  $i$  in our enumeration.  
 2. aa Is aaaaaa in the enumeration?  $\rightarrow$   
 3. aaa

Claim:  $L((a+b)^*)$  is countably infinite.

Proof: attempt: alpha order 2nd attempt. Shortlex order.

- first ordering principle  
length
- within a length, alpha order.

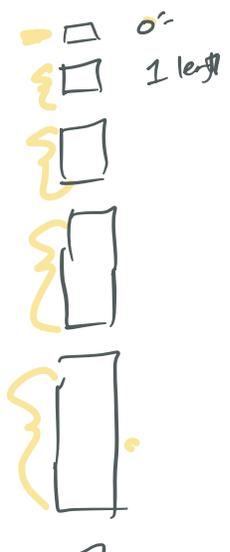
Problem  $\exists$  an infinite # of strings before  $b$  in the "enumeration".

This enumeration scheme works because for each string,  $\exists$  a finite # of strings that precede it in the enumeration.

see proof later.  $\square$

1.  $\epsilon$
2. a
3. b.
4. aa
5. ab
6. ba
7. bb
- ...

abbab



Claim:  $\mathbb{Q}^+$ , the set of positive rationals, is countably infinite.



Proof: Clearly it is infinite, since it contains  $\mathbb{N}$ .  $\mathbb{N} \subseteq \mathbb{Q}^+$

How can we enumerate the rationals?

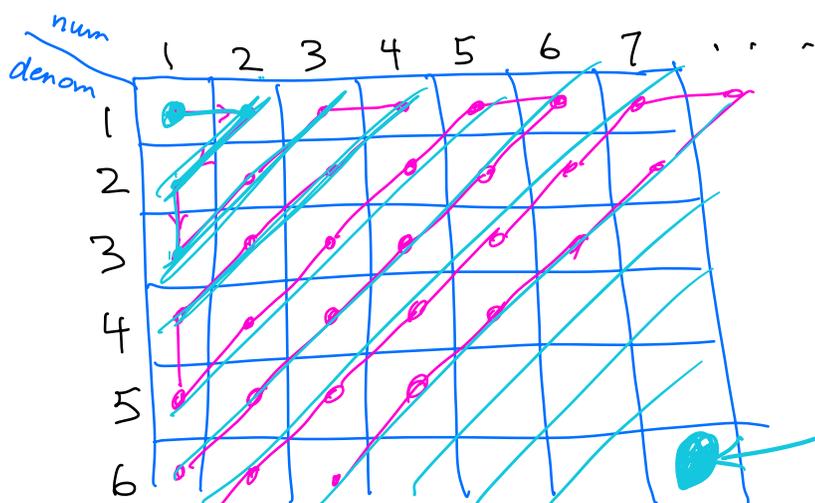
1<sup>st</sup> attempt (doomed)

1.  $\frac{1}{1}$
2.  $\frac{1}{2}$
3.  $\frac{1}{3}$
4.  $\frac{1}{4}$
- ...

Problem: never get to  $\frac{2}{3}$ .

2<sup>nd</sup> attempt - in order of size - but what is the smallest positive rational?

3<sup>rd</sup> attempt - A pos rational is a pos int numerator and a pos int denominator.



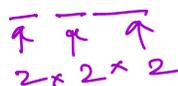
1.  $\frac{1}{1}$  ✓
2.  $\frac{2}{1}$  ✓
3.  $\frac{1}{2}$  ✓
4.  $\frac{1}{3}$  ✓
- ...
- ~~5.  $\frac{2}{2}$~~
- ~~6.  $\frac{1}{2}$~~
- ~~7.  $\frac{3}{2}$~~
- ~~8.  $\frac{2}{2}$~~

This is not a diagonalization proof -  
 just a nice method for enumerating the cross product  
 of two countably infinite sets.

Going back to our shortlex enumeration.. of strings over  $\Sigma$ .

Claim:  $\forall$  length  $i$ ,  $\exists |\Sigma|^i$  strings of that length.

$i = 3, \Sigma = \{a, b\}$



Shortlex:

1.  $\epsilon$
2. a
3. b
- ...

Claim:  $\forall$  strings  $w$   $\exists$  a finite number of strings  
 over  $\Sigma$  that precede it in shortlex order.

Proof: Let  $|w| = i$ .

$\exists$  a finite # of lengths  $j$  such that  $j < i$   
 Each length has a finite # of strings of that  
 length.

Also,  $\exists$  a finite # of strings of same length  
 as  $w$ .

Note: Shortlex works on any alphabet

If  $\nexists$  an implicit ordering on the symbols of the  
 alphabet, you can impose one arbitrarily.

eg  $\{\Delta, \circ, \square\} \quad \Delta < \circ < \square$

$\{0, 1\} \quad 0 < 1$

Is every set countable?

How about the set of infinite bit strings?

0 1 0 0 0 1 0 0 1 1 1 1 0 1 0 1 1 1 0 1 0 1 0 0 0 ...

Madelaine says "YES", she has an enumeration.

I say: I can construct an infinite bitstring that is NOT in your enumeration.

Madelaine's enumeration:

T	1	2	3	4	...	...	...	
1	0	1	0	1	1	0	0	1
2	1	0	1	0	0	1	1	1
3	0	0	0	0	0	0	0	0
4	1	1	1	1	0	1	0	1
5	1	1	1	0	1	0	1	1
6	0	1	1	0	0	0	1	0
7	1	0	1	0	0	1	0	1

$\bar{D} = 1110011$

$D$  = the diagonal bitstring.

$\rightarrow D[i] = T[i][i]$

$\rightarrow \bar{D}$  is the "flip" of the diagonal ie

$$\bar{D}[i] = 1 - D[i] = 1 - T[i][i]$$

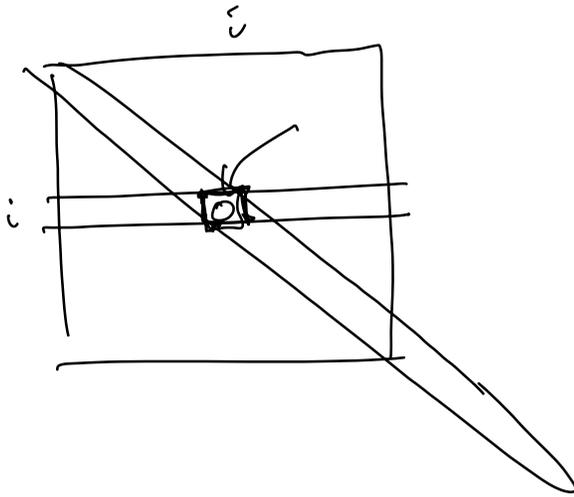
Claim:  $\bar{D}$  is not enumerated.

Proof: If  $\bar{D}$  is enumerated, then it has a rank

$$r.$$

Then  $\bar{D}[r] = \underbrace{T[r][r]}_{\bar{D}}$

But  $\bar{D}[r] = 1 - T[r][r], \Rightarrow \Leftarrow$



Georg Cantor

$$|\mathbb{N}| = \aleph_0$$

Aleph nought

- ∴ No enumeration can exist for infinite bitstrings.
- ∴ inf bitstrings are uncountably infinite.

$\mathbb{R}$  = the set of real numbers.  $\mathbb{R}_{[0,1)} = \mathbb{R} \cap [0...1)$

Claim:  $\mathbb{R}_{[0,1)}$  is not countable.

Proof: BWOC.  $\S$   $\mathbb{R}_{[0,1)}$  is countable. Then  $\exists$  an enumeration:

1	.0 <sup>1</sup>	1	9	0	1	3	4	8	7	6	...
2	.9	2 <sup>3</sup>	1	2	2	3	7	9	1	0	...
3	.3	9	2 <sup>3</sup>	2	2	9	3	4	1	6	...
4	.2	2	2	2 <sup>3</sup>	0	0	0	0	0	0	...
5	.5	0	0	0	0 <sup>1</sup>	0	0	0	0	0	...
6	.7	6	0	1	9	2 <sup>3</sup>	8	7	8	7	...
7	.0	0	0	1	4	1	6 <sup>7</sup>	2	6	6	...
8	.1	0	0	1	9	2	1	9 <sup>0</sup>	3	4	...
...											...

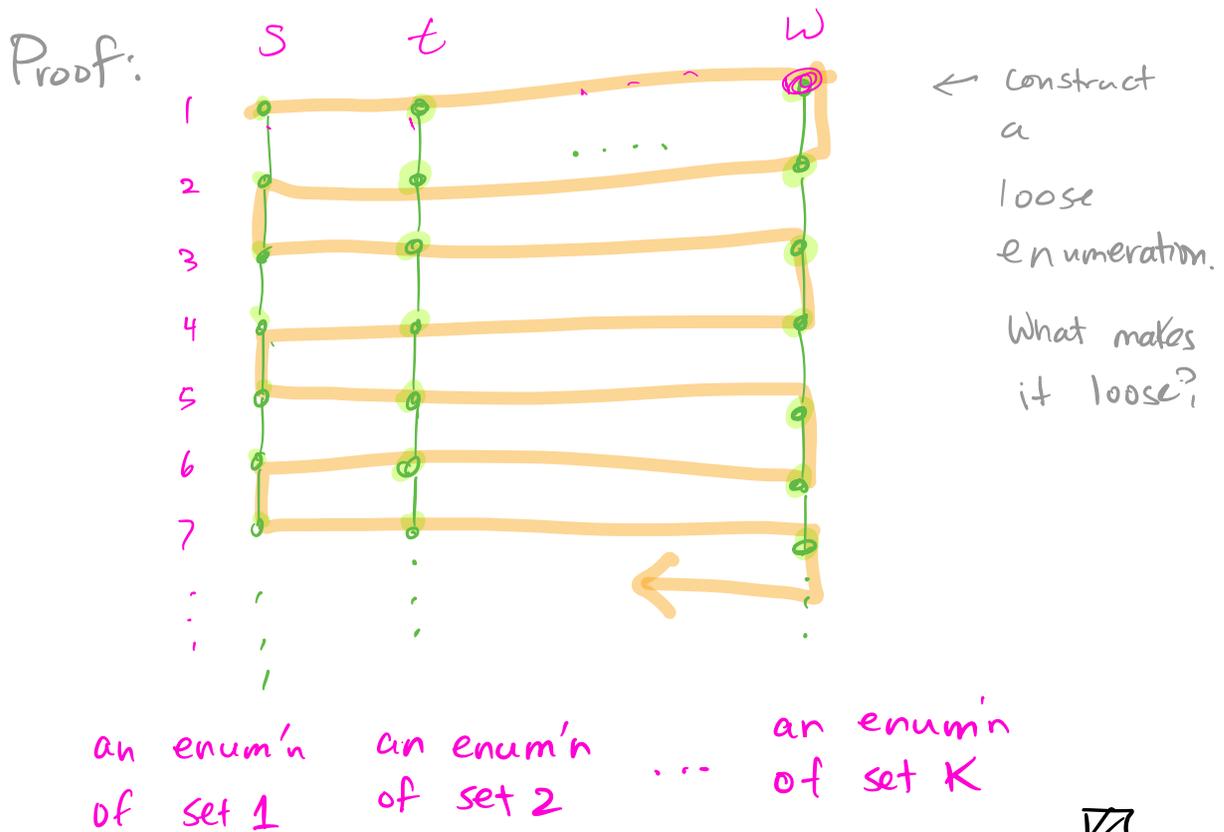
Then the number  $0.13331370\dots$ ,  
 constructed by taking each digit along the diagonal  
 and adding 1 to it (mod 10) is not in the

enumeration!

⇒ it is not an enumeration.

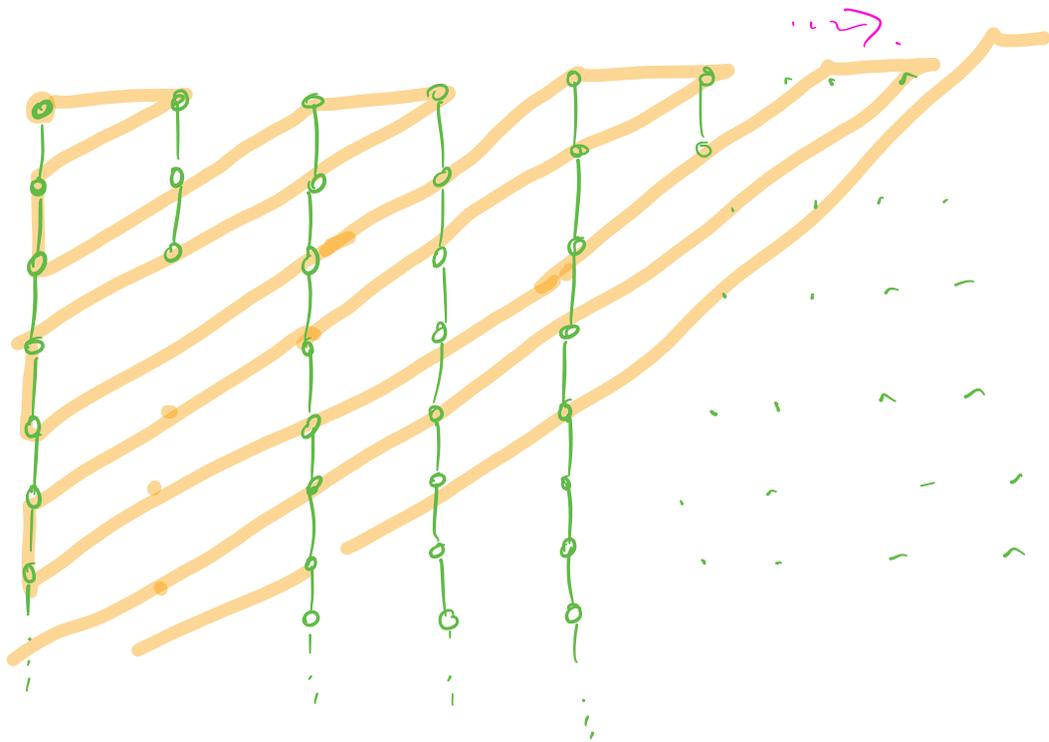
Furthermore,  $\forall$  proposed enumerations,  
you can always construct a real number  
that is not enumerated.  $\square$

Claim: The union of  $k$  countably infinite  
sets is also countably infinite.



Claim: The union of a countably infinite number of countable sets is countable.

Proof:



It is "loose" because:

- Some of the sets may "end early" (be finite).

But a loose enumeration is OK.



Warning: If a problem (on assignment or test) asks you to show something is countably infinite by giving an enum. scheme, I want to see the enumeration scheme, not the application of one of the theorems above.